

Taskflow: Programming System for building High-performance EDA Applications

How can we make it easier to program heterogeneous EDA applications?

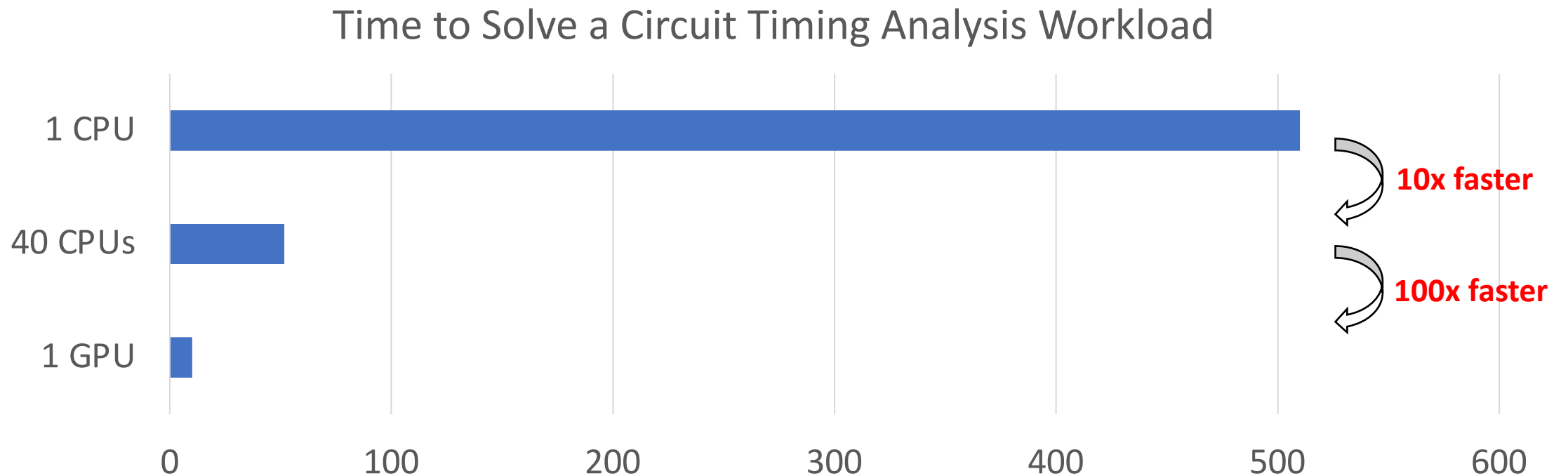
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Department of Electrical and Computer Engineering
University of Utah, Salt Lake City, UT

<https://tsung-wei-huang.github.io/>



Why Parallel Computing?

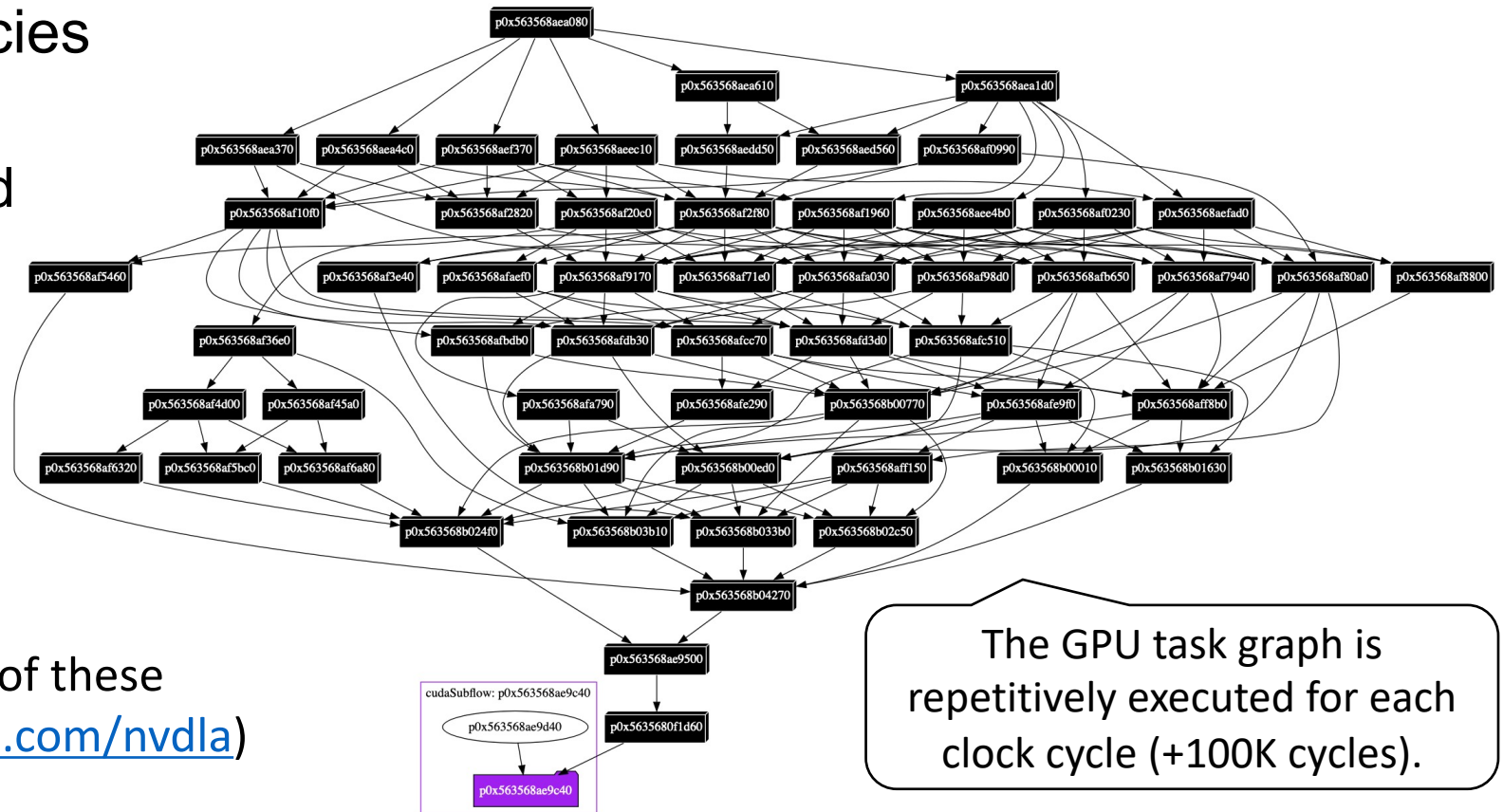
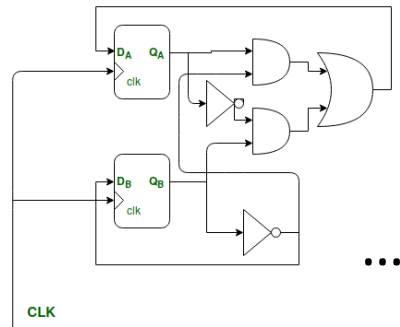
- It's critical to advance your application performance



Today's Workload is Very Complex

- GPU-accelerated circuit analysis on a design of 500M gates

- >100 kernels
- >100 dependencies
- >500s to finish
- >10hrs turnaround



The GPU task graph is repetitively executed for each clock cycle (+100K cycles).

What are the output values of these 500M gates? (<https://github.com/nvdla>)

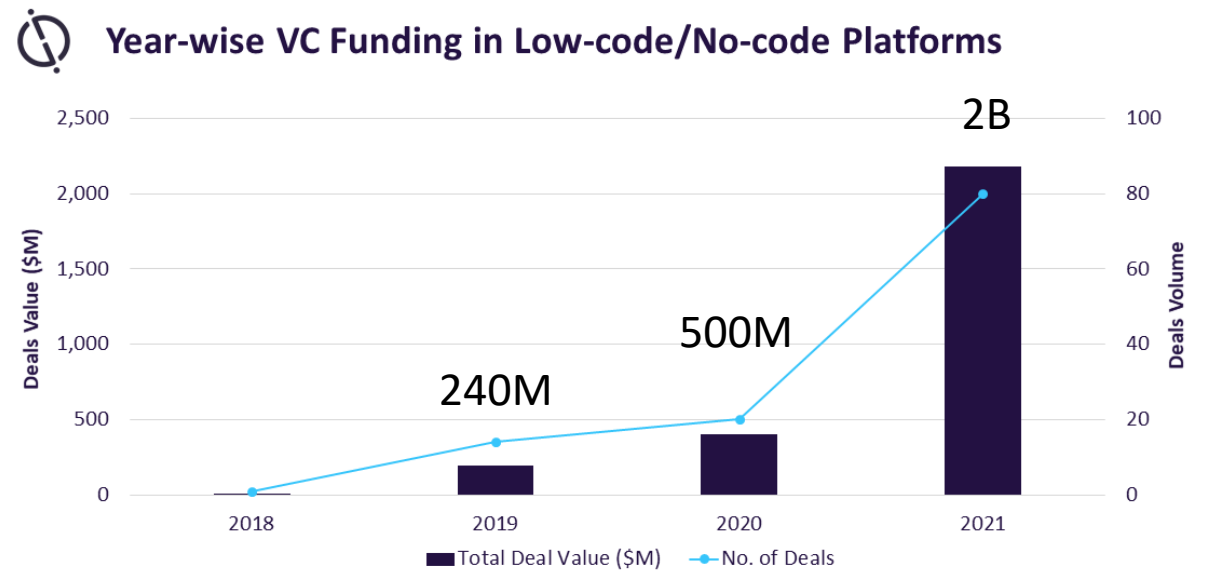
Parallel Programming is VERY Challenging

- You need to deal with A LOT OF technical details
 - Parallelism abstraction (task, data, concurrent data structures, etc.)
 - Concurrency control
 - Task/data race avoidance
 - Dependency constraints
 - Load balancing
 - Scheduling efficiencies
 - ...



I want to focus more on my applications ...

How can we make scientific software researchers' (your) lives easier?



Source: GlobalData Disruptor Intelligence Center – Deals Database

GlobalData.



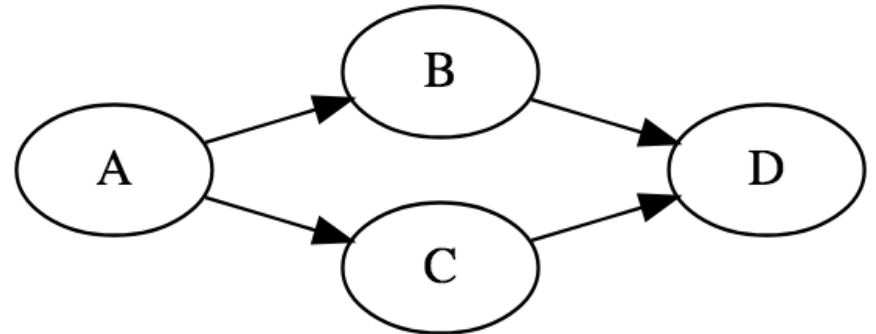
Taskflow offers a solution

*How can we make it easier for C++ developers to quickly write parallel and heterogeneous programs with **high performance scalability** and **simultaneous high productivity**?*



“Hello World” in Taskflow

```
#include <taskflow/taskflow.hpp> // Taskflow is header-only
int main(){
    tf::Taskflow taskflow;
    tf::Executor executor;
    auto [A, B, C, D] = taskflow.emplace(
        [] () { std::cout << "TaskA\n"; },
        [] () { std::cout << "TaskB\n"; },
        [] () { std::cout << "TaskC\n"; },
        [] () { std::cout << "TaskD\n"; }
    );
    A.precede(B, C); // A runs before B and C
    D.succeed(B, C); // D runs after B and C
    executor.run(taskflow).wait(); // submit the taskflow to the executor
    return 0;
}
```



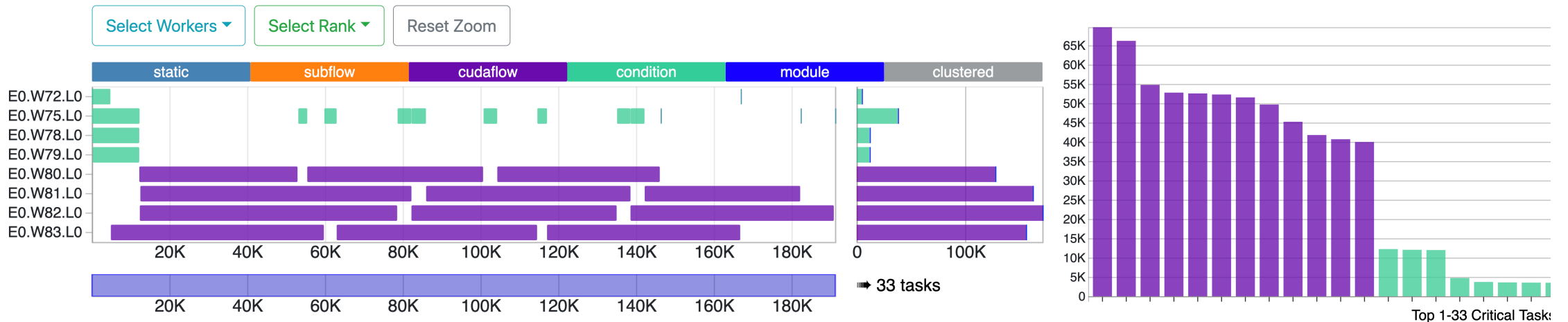
Drop-in Integration

- Taskflow is header-only – *no wrangle with installation*

```
~$ git clone https://github.com/taskflow/taskflow.git # clone it only once
~$ g++ -std=c++17 simple.cpp -I taskflow/taskflow -O2 -pthread -o simple
~$ ./simple
TaskA
TaskC
TaskB
TaskD
```

Built-in Profiler/Visualizer

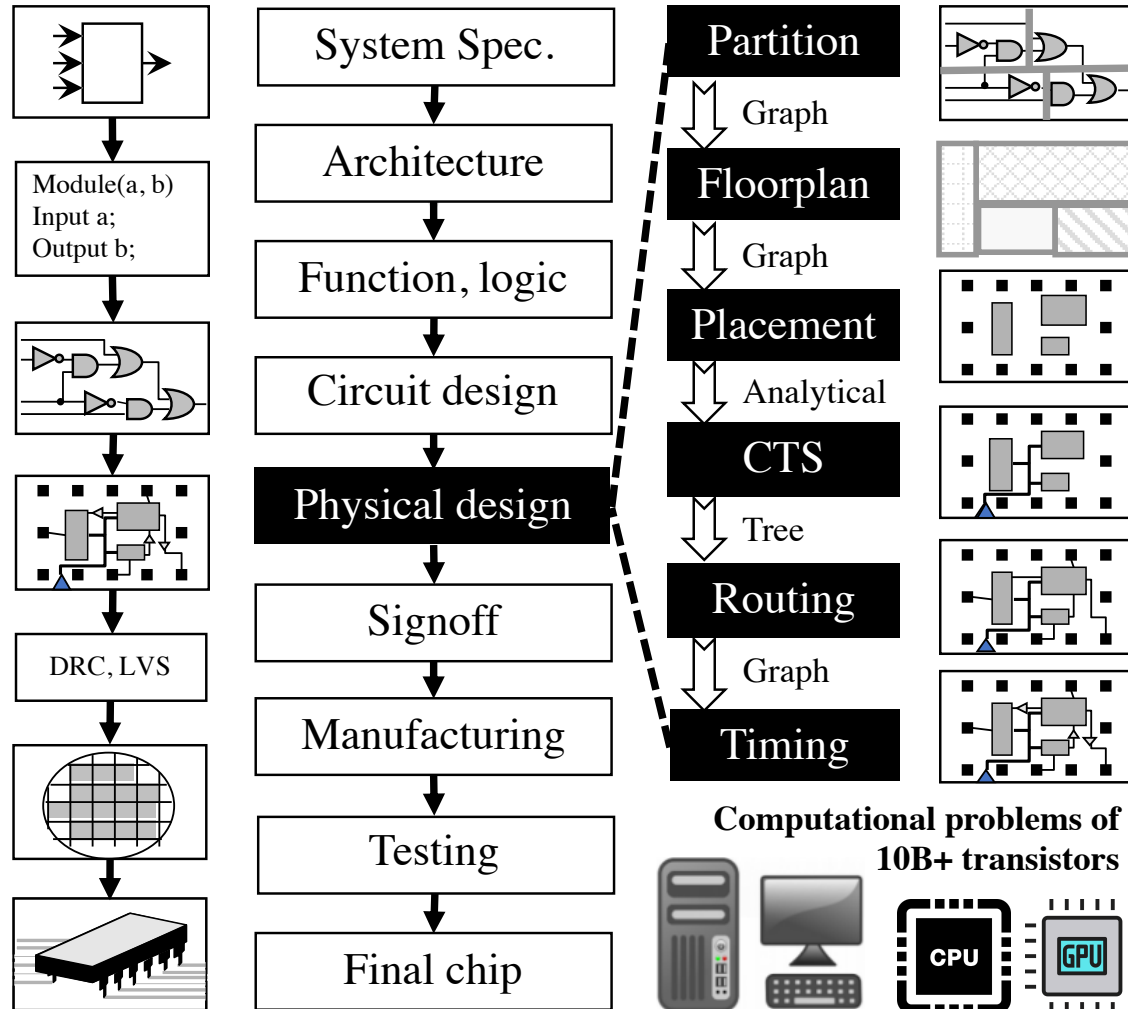
```
# run the program with the environment variable TF_ENABLE_PROFILER enabled
~$ TF_ENABLE_PROFILER=simple.json ./simple
~$ cat simple.json
[
{"executor": "0", "data": [{"worker": 0, "level": 0, "data": [{"span": [172, 186], "name"}]}]}]
# paste the profiling json data to https://taskflow.github.io/tfprof/
```



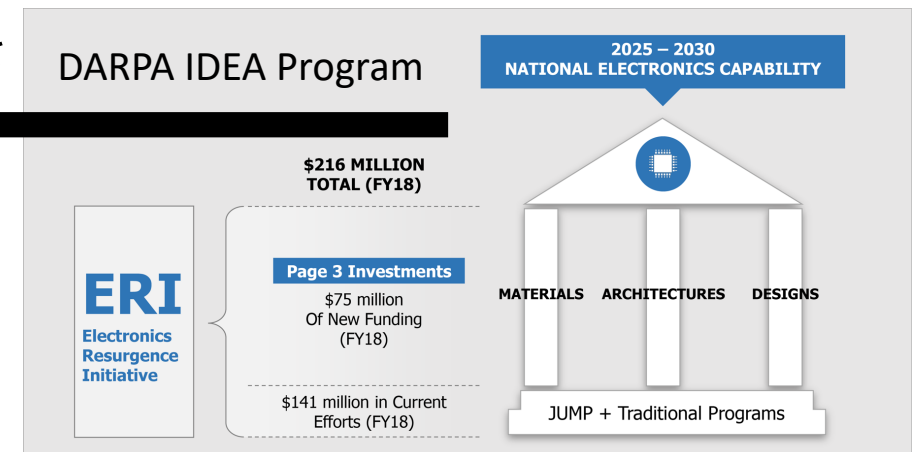
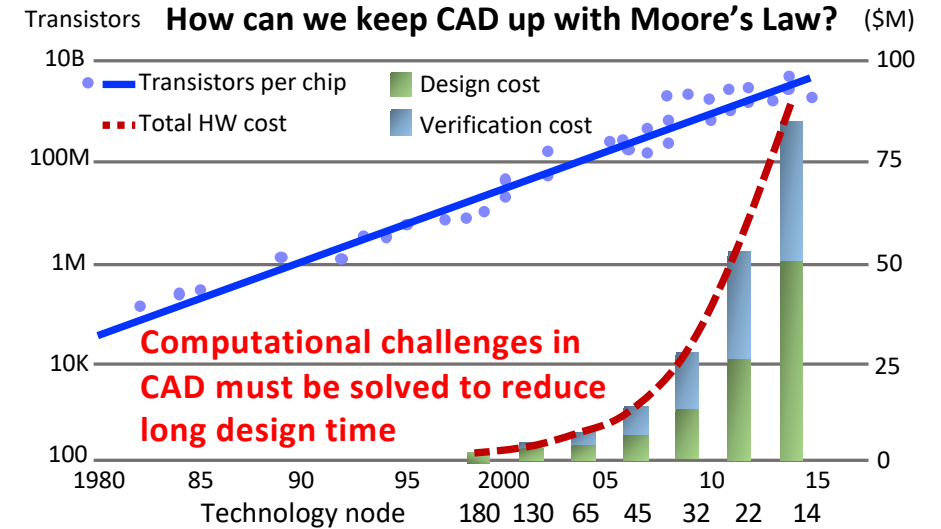
Agenda

- **Express your parallelism in the right way**
- Parallelize your applications using Taskflow
- Boost performance in CAD applications

Motivation: Parallelizing CAD/EDA Tools



24-hour design cycle



We Invested a lot in Existing Tools ...

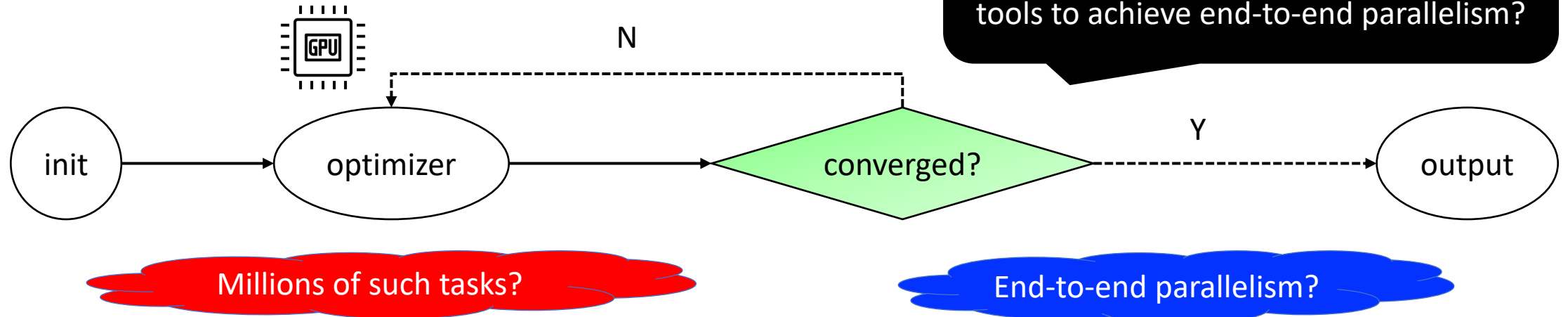


Two Big Problems of Existing Tools

- EDA has **complex task dependencies**
 - **Example**: analysis algorithms compute the circuit network of multi-millions of nodes and dependencies
 - **Problem**: existing tools are often good at loop parallelism but weak in expressing heterogeneous task graphs at this large scale
- EDA has **complex control flow**
 - **Example**: synthesis algorithms make essential use of *dynamic control flow* to implement various patterns
 - Combinatorial optimization (e.g., graph algorithms, discrete math)
 - analytical methods (e.g., physical synthesis)
 - **Problem**: existing tools are *directed acyclic graph* (DAG)-based and do not anticipate control flow in the graph, lacking *end-to-end* parallelism

Example: An Iterative Optimizer

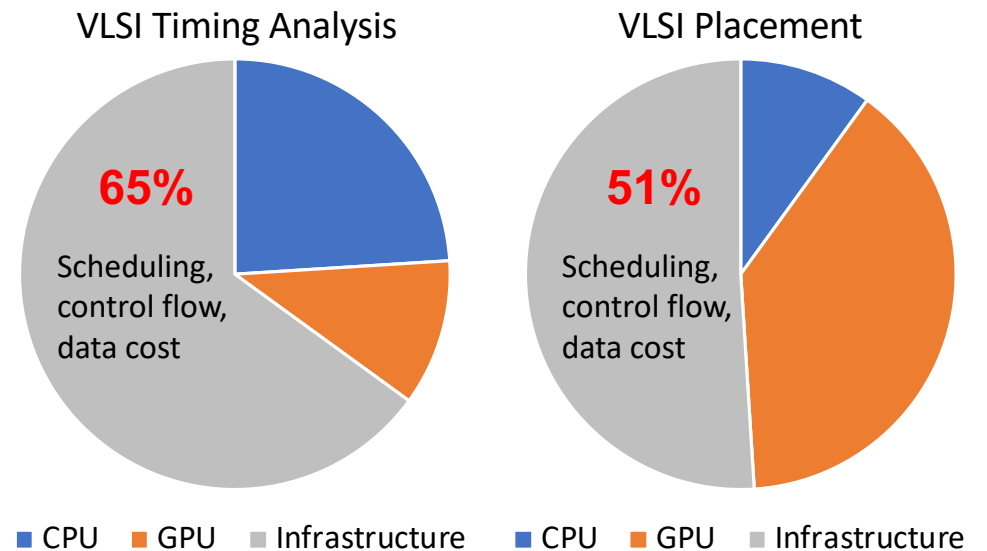
- Four computational tasks with dynamic control flow
 - #1: starts with `init` task
 - #2: enters the `optimizer` task (e.g., GPU math solver)
 - #3: checks if the optimization converged
 - No: loops back to `optimizer`
 - Yes: proceeds to `stop`
 - #4: outputs the result



Need a New Parallel Programming System

While designing parallel algorithms is non-trivial, what makes parallel programming an enormous challenge is the **infrastructure work** of “*how to efficiently express dependent tasks along with algorithmic control flow and schedule them across heterogeneous computing resources*”

- **VLSI timing analysis (ICCAD'20)**
 - up to **65%** runtime on infrastructure
 - 24% on CPU and 11% on GPU
- **VLSI placement (TCAD'21)**
 - up to **51%** runtime on infrastructure
 - 10% on CPU and 39% on GPU



Agenda

- Express your parallelism in the right way
- **Parallelize your applications using Taskflow**
- Boost performance in CAD applications

Revisit “Hello World” in Taskflow (TPDS’22)

```
#include <taskflow/taskflow.hpp> // Taskflow is header-only
int main(){
    tf::Taskflow taskflow;
    tf::Executor executor;
    auto [A, B, C, D] = taskflow.emplace(
        [] () { std::cout << "TaskA\n"; },
        [] () { std::cout << "TaskB\n"; },
        [] () { std::cout << "TaskC\n"; },
        [] () { std::cout << "TaskD\n"; }
    );
    A.precede(B, C); // A runs before B and C
    D.succeed(B, C); // D runs after B and C
    executor.run(taskflow).wait(); // submit the taskflow to the executor
    return 0;
}
```

A new **control taskflow graph (CTFG)** programming model

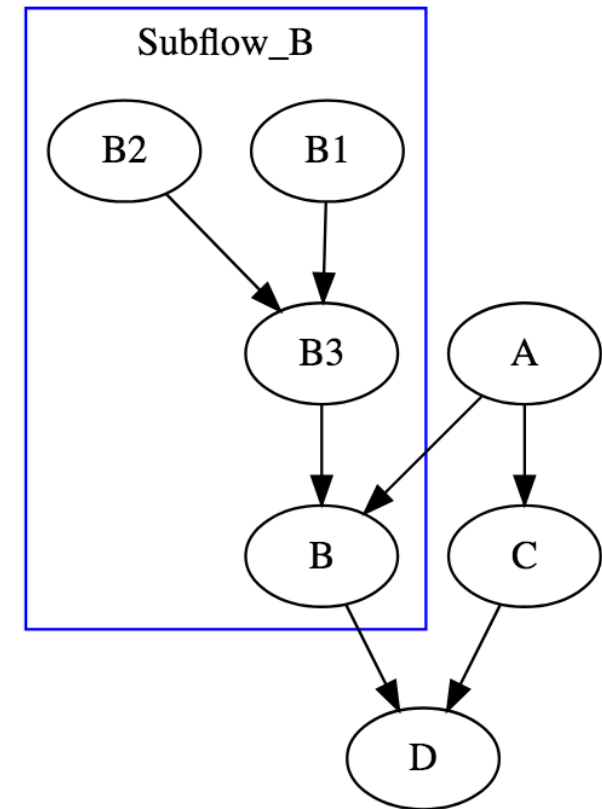
1. Static tasking
 2. Dynamic tasking
 3. Conditional tasking
 4. Heterogeneous tasking
 5. Pipeline tasking
- ... (more on <https://taskflow.github.io/>)

#2: Dynamic Tasking (Subflow)

```
// create three regular tasks  
tf::Task A = tf.emplace([](){}).name("A");  
tf::Task C = tf.emplace([](){}).name("C");  
tf::Task D = tf.emplace([](){}).name("D");
```

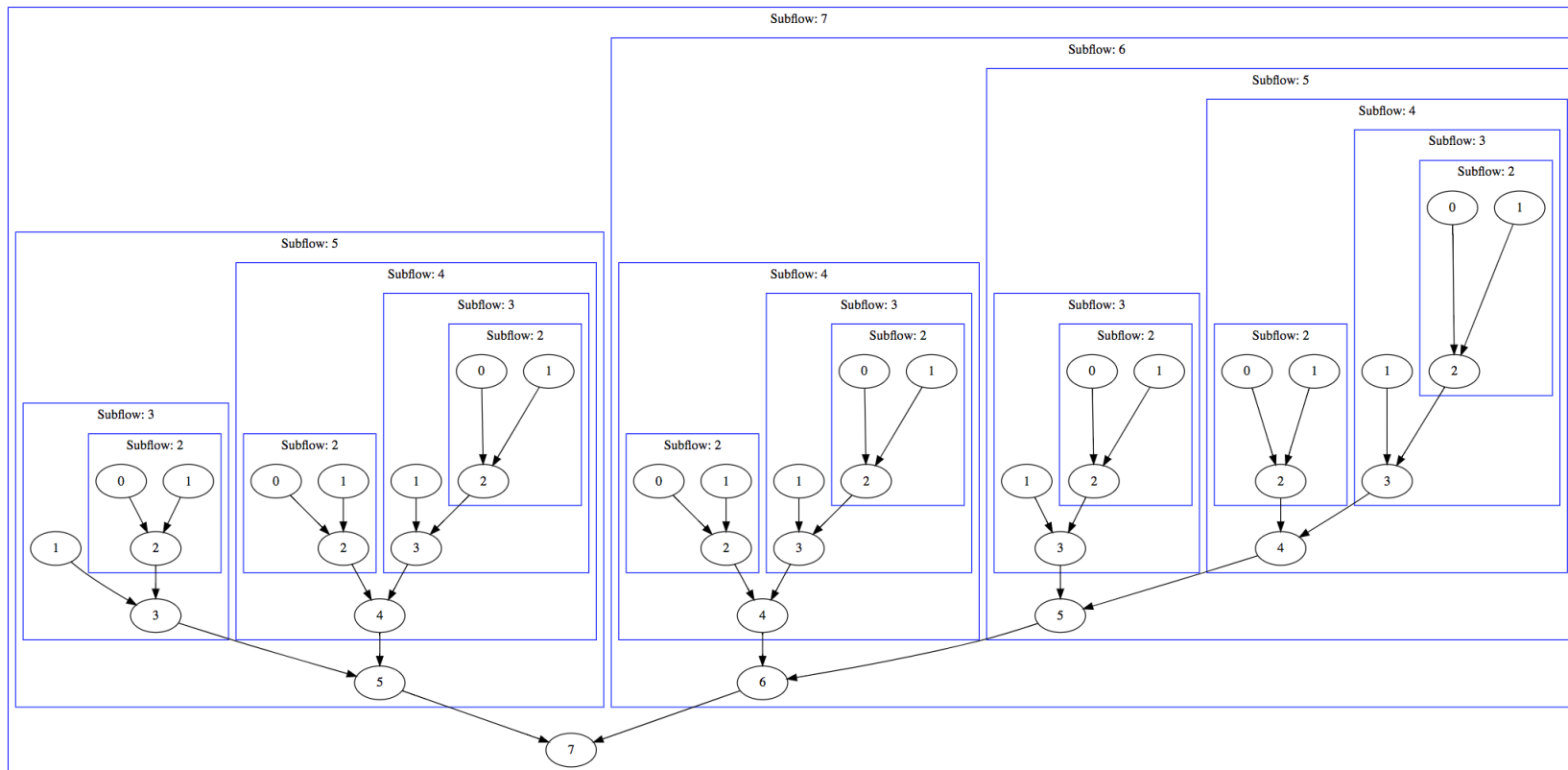
```
// create a subflow graph (dynamic tasking)  
tf::Task B = tf.emplace([] (tf::Subflow& subflow) {  
    tf::Task B1 = subflow.emplace([](){}).name("B1");  
    tf::Task B2 = subflow.emplace([](){}).name("B2");  
    tf::Task B3 = subflow.emplace([](){}).name("B3");  
    B1.precede(B3);  
    B2.precede(B3);  
}).name("B");
```

```
A.precede(B); // B runs after A  
A.precede(C); // C runs after A  
B.precede(D); // D runs after B  
C.precede(D); // D runs after C
```



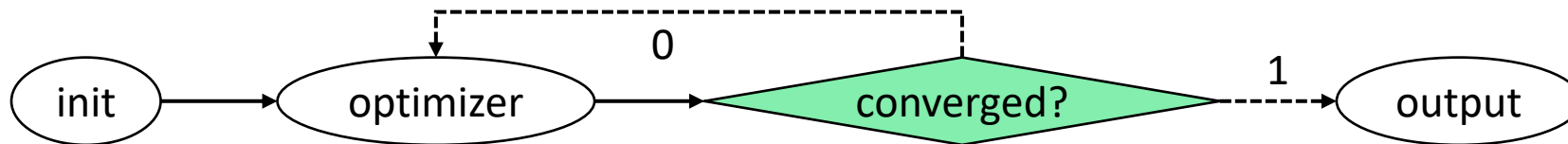
Subflow can be Nested and Recursive

- Find the 7th Fibonacci number using subflow
 - $\text{Fib}(n) = \text{Fib}(n-1) + \text{Fib}(n-2)$



#3: Conditional Tasking (if-else)

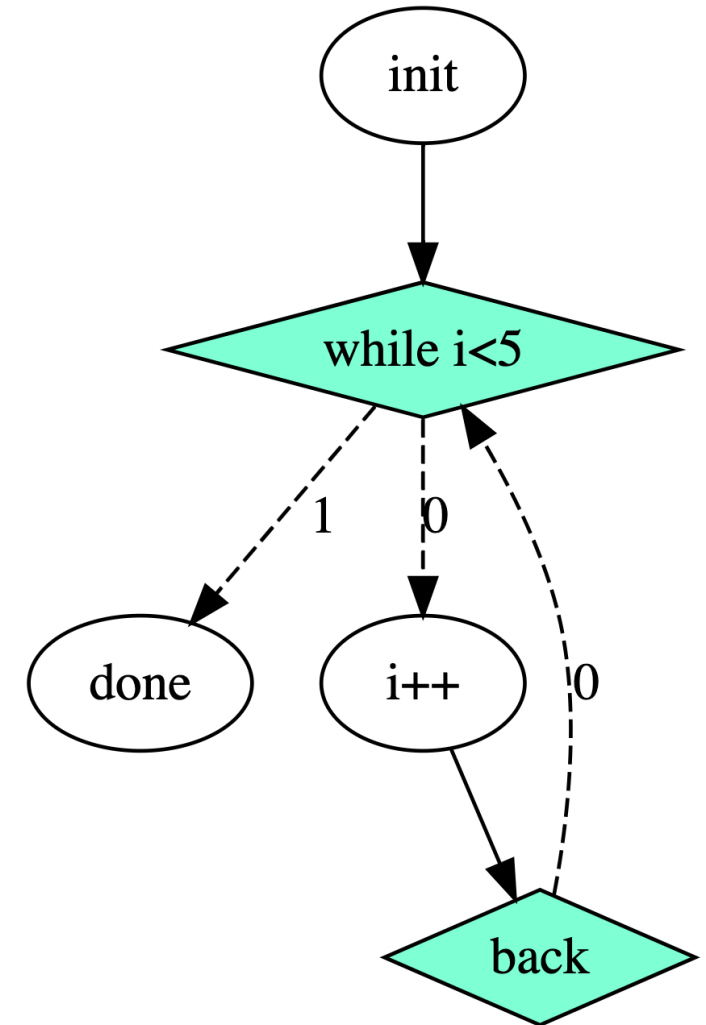
```
auto init          = taskflow.emplace([&]() { initialize_data_structure(); } )  
                  .name("init");  
auto optimizer     = taskflow.emplace([&]() { matrix_solver(); } )  
                  .name("optimizer");  
auto converged     = taskflow.emplace([&]() { return converged() ? 1 : 0 ; } )  
                  .name("converged");  
auto output       = taskflow.emplace([&]() { std::cout << "done!\n"; } );  
                  .name("output");  
  
init.precede(optimizer);  
optimizer.precede(converged);  
converged.precede(optimizer, output); // return 0 to the optimizer again
```



*Condition task enables in-graph control flow to achieve **end-to-end** parallelism*

#3: Conditional Tasking (iterative loop)

```
tf::Taskflow taskflow;  
int i;  
auto [init, cond, body, back, done] = taskflow.emplace(  
    [&]() { std::cout << "i=0"; i=0; },  
    [&]() { std::cout << "while i<5\n"; return i < 5 ? 0 : 1; },  
    [&]() { std::cout << "i++=" << i++ << "\n"; },  
    [&]() { std::cout << "back\n"; return 0; },  
    [&]() { std::cout << "done\n"; }  
);  
init.precede(cond);  
cond.precede(body, done);  
body.precede(back);  
back.precede(cond);
```

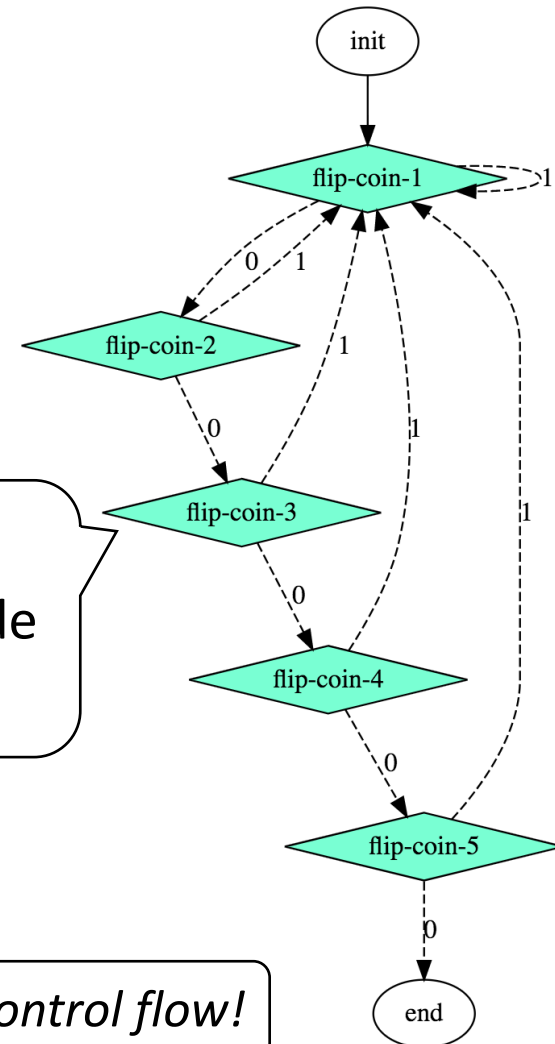


#3: Conditional Tasking (random loops)

```
auto A = taskflow.emplace([&](){});  
auto B = taskflow.emplace([&]() { return rand()%2; } );  
auto C = taskflow.emplace([&]() { return rand()%2; } );  
auto D = taskflow.emplace([&]() { return rand()%2; } );  
auto E = taskflow.emplace([&]() { return rand()%2; } );  
auto F = taskflow.emplace([&]() { return rand()%2; } );  
auto G = taskflow.emplace([&](){});  
A.precede(B).name("init");  
B.precede(C, B).name("flip-coin-1");  
C.precede(D, B).name("flip-coin-2");  
D.precede(E, B).name("flip-coin-3");  
E.precede(F, B).name("flip-coin-4");  
F.precede(G, B).name("flip-coin-5");  
G.name("end");
```

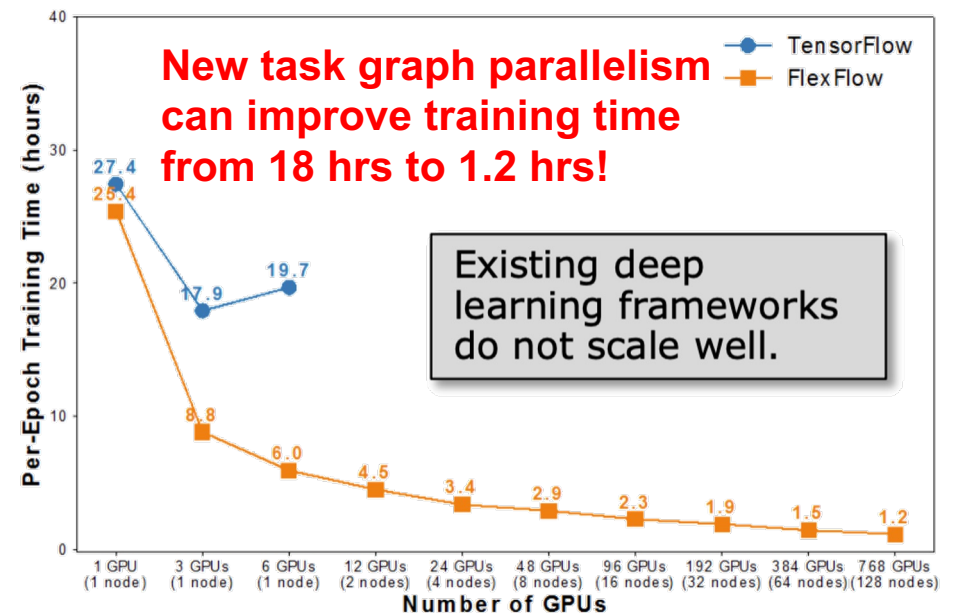
Each task flips a binary coin to decide the next path

You can describe non-deterministic, nested control flow!



Existing Frameworks on Control Flow?

- **Expand a task graph across fixed-length iterations**
 - Large graph size linearly proportional to decision points
- **Unknown or non-deterministic iterations?**
 - Expensive dynamic tasks executing “if-else” on the fly
- **Dynamic control-flow tasks?**
 - Client-side partition
- **Same problem in large-scale ML**
 - TensorFlow with RNN (EuroSys’18)
 - FlexFlow (MLSys’19, ICML’18)
 - DGL (CoRR’19)
 - DOE 2022 funding preview (Dr. Finkel)



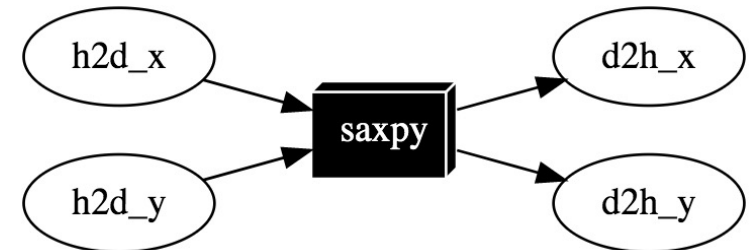
#4: Heterogeneous Tasking

```
const unsigned N = 1<<20;
std::vector<float> hx(N, 1.0f), hy(N, 2.0f);
float *dx{nullptr}, *dy{nullptr};
auto allocate_x = taskflow.emplace([&]() { cudaMalloc(&dx, 4*N); });
auto allocate_y = taskflow.emplace([&]() { cudaMalloc(&dy, 4*N); });
```

```
auto cudaflow = taskflow.emplace([&](tf::cudaFlow& cf) {
    auto h2d_x = cf.copy(dx, hx.data(), N); // CPU-GPU data transfer
    auto h2d_y = cf.copy(dy, hy.data(), N);
    auto d2h_x = cf.copy(hx.data(), dx, N); // GPU-CPU data transfer
    auto d2h_y = cf.copy(hy.data(), dy, N);
    auto kernel = cf.kernel((N+255)/256, 256, 0, saxpy, N, 2.0f, dx, dy);
    kernel.succeed(h2d_x, h2d_y).precede(d2h_x, d2h_y);
});
```

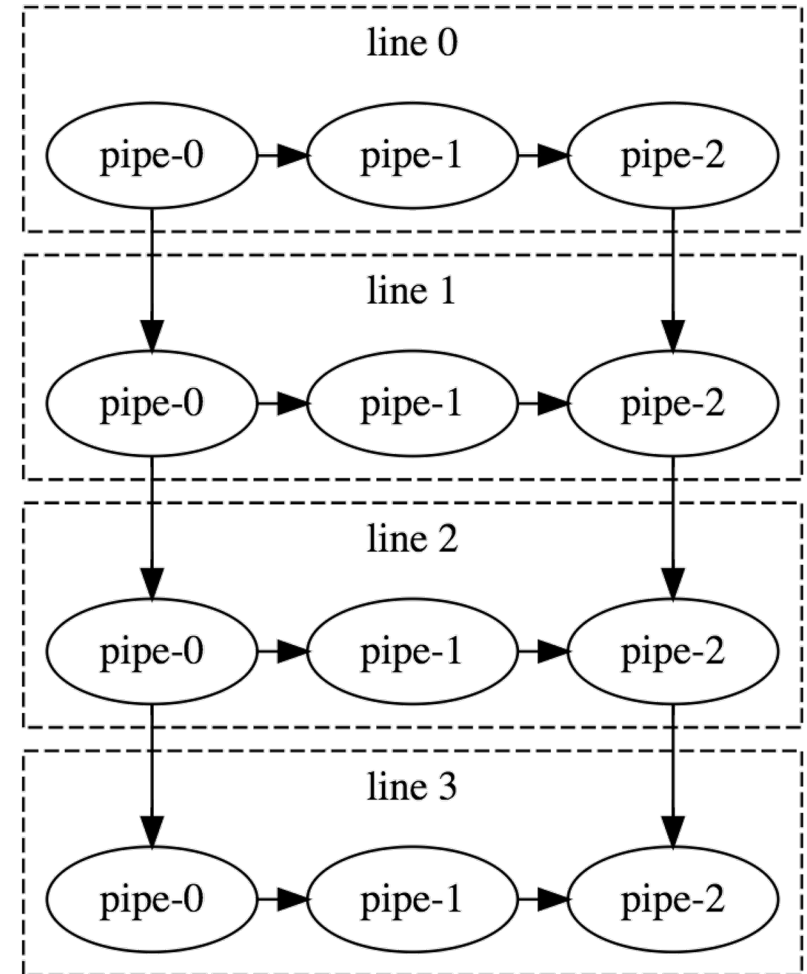
```
cudaflow.succeed(allocate_x, allocate_y);
executor.run(taskflow).wait();
```

cudaFlow automatically transforms an application GPU task graph to an optimized “CUDA graph”



5: Pipeline Tasking (HPDC'22)

```
std::array<int, 4> buffer;  
tf::Pipeline pl(4,  
  tf::Pipe {tf::PipeType::SERIAL, [&buffer](tf::Pipeflow & pf) {  
    if (pf.token() == 5) {  
      pf.stop();  
      return;  
    }  
    buffer[pf.line()] = pf.token();  
  }},  
  tf::Pipe {tf::PipeType::PARALLEL, [&buffer](tf::Pipeflow & pf) {  
    buffer[pf.line()] = buffer[pf.line()] + 1;  
  }},  
  tf::Pipe {tf::PipeType::SERIAL, [&buffer](tf::Pipeflow & pf) {  
    buffer[pf.line()] = buffer[pf.line()] + 1;  
  }}  
);  
auto task = taskflow.composed_of(pl);  
executor.run(taskflow).wait();
```



Submit Taskflow to Executor

- Executor manages a set of threads to run taskflows
 - All execution methods are *non-blocking*
 - All execution methods are *thread-safe*

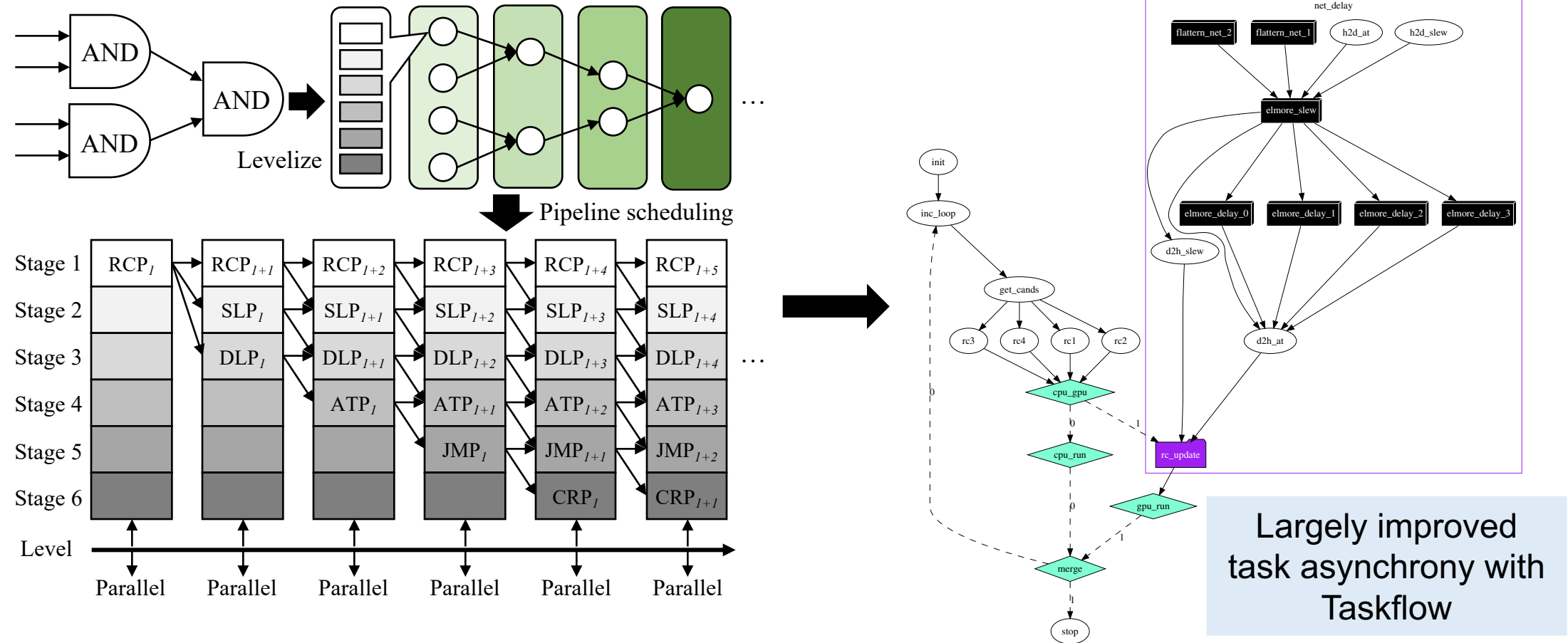
```
{  
  tf::Taskflow taskflow1, taskflow2, taskflow3;  
  tf::Executor executor;  
  // create tasks and dependencies  
  // ...  
  auto future1 = executor.run(taskflow1);  
  auto future2 = executor.run_n(taskflow2, 1000);  
  auto future3 = executor.run_until(taskflow3, [i=0]() { return i++>5 });  
  executor.wait_for_all(); // wait for all the above tasks to finish  
}
```

Agenda

- Express your parallelism in the right way
- Parallelize your applications using Taskflow
- **Boost performance in CAD applications**

Case Study 1: Timing Analysis (TCAD'21)

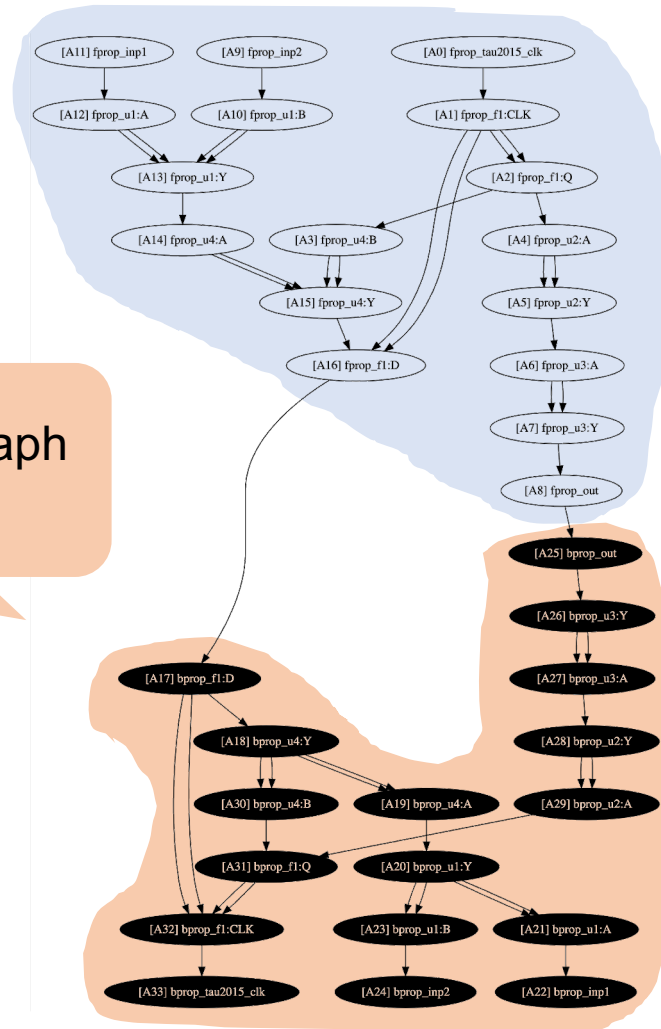
- Analyzing the timing of large circuits can take *hours* to finish



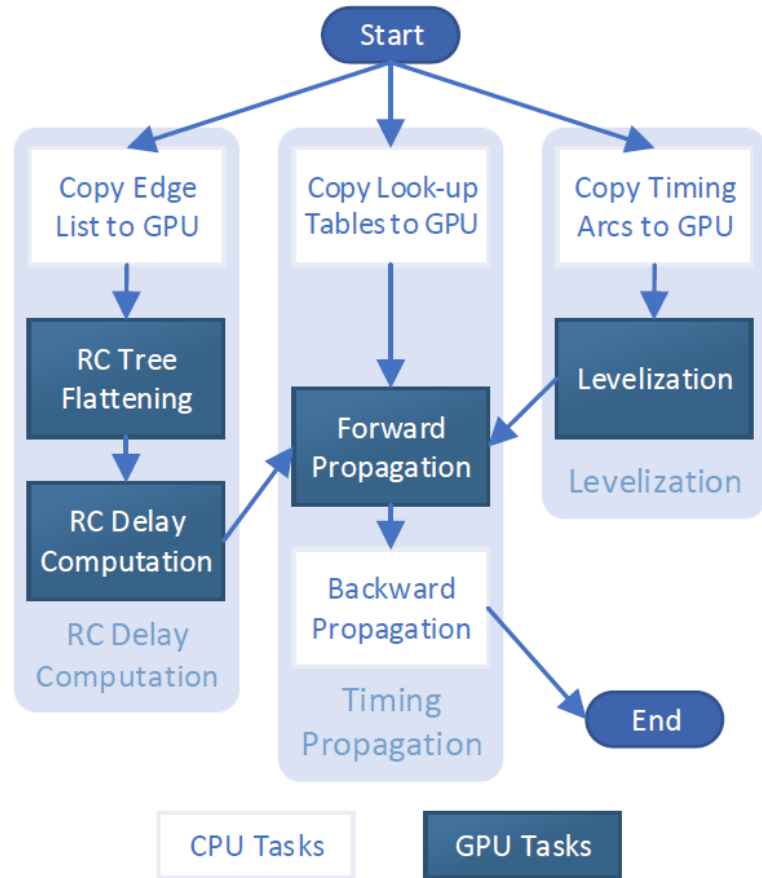
Case Study: Timing Analysis (cont'd)

Forward propagation task graph
(RC, arrival time, slew, etc.)

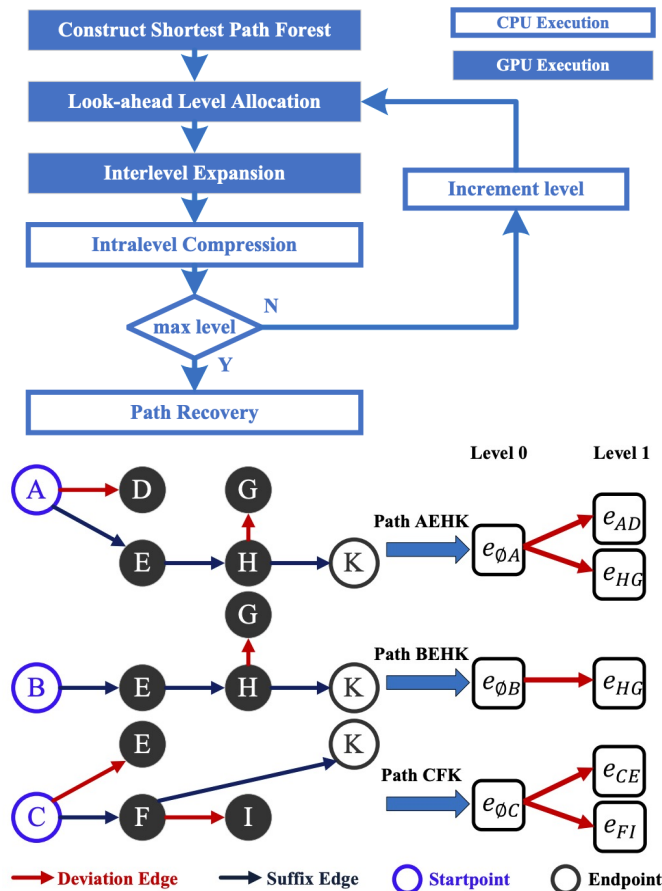
Backward propagation task graph
(constraints, slack, etc.)



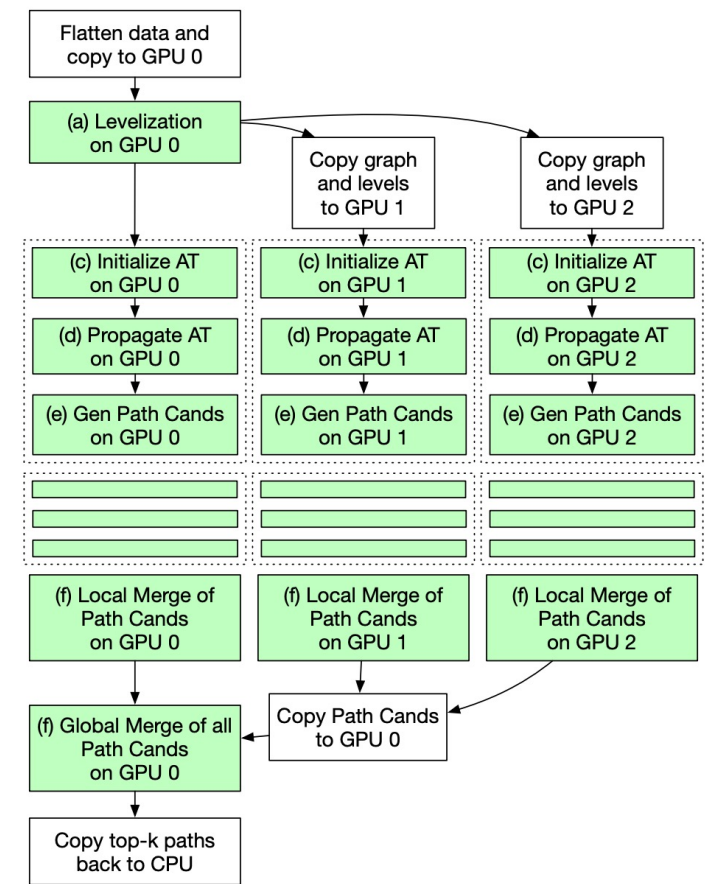
Case Study 1: Timing Analysis (cont'd)



GPU-based graph analysis (ICCAD'20)



GPU-based path analysis (DAC'21)



GPU-based CPPR (ICCAD'21)

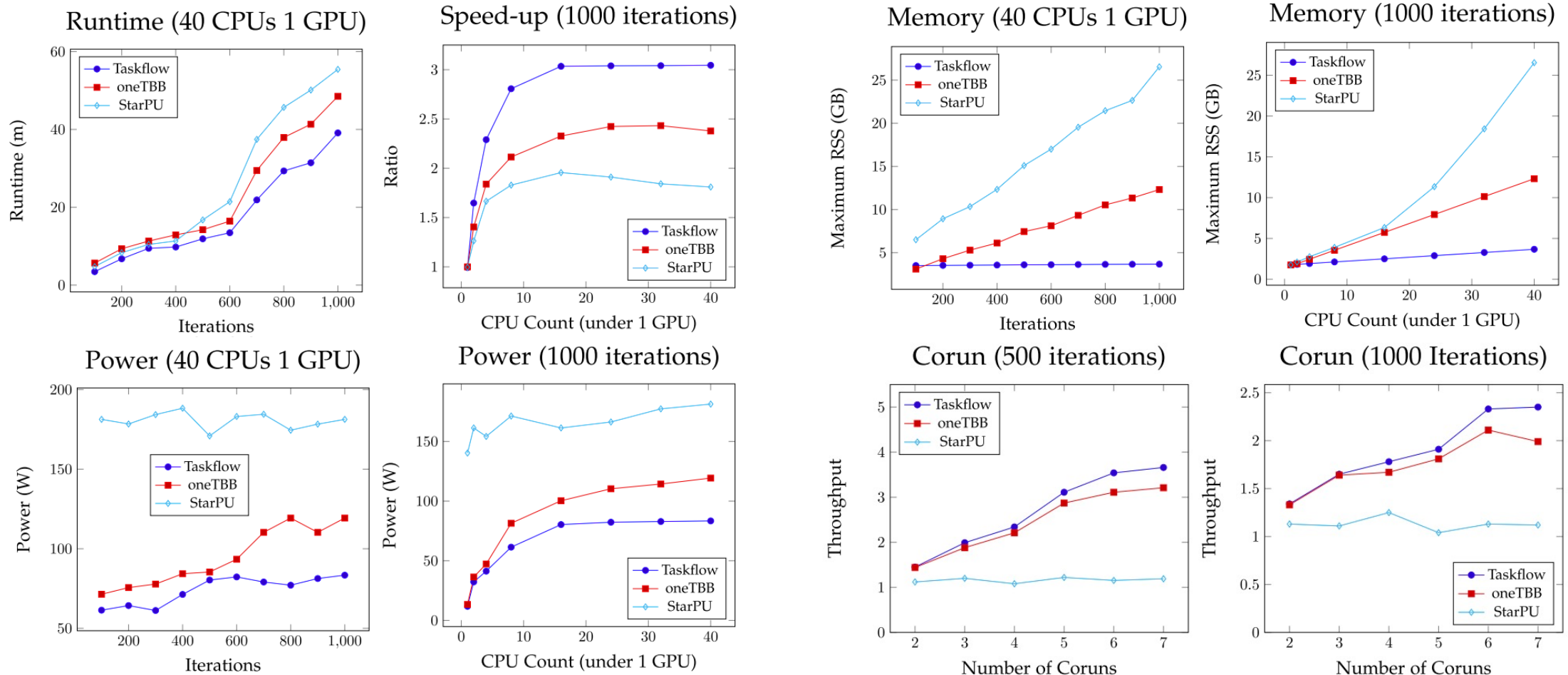
Case Study 1: Timing Analysis (cont'd)

- **Path-based timing analysis (DAC'21, Best Paper in TAU'21)**
 - leon3mp (1.6M gates): **611x speed-up** over 1 CPU (**44x** over 40 CPUs)
 - netcard (1.5M gates): **367x speed-up** over 1 CPU (**46x** over 40 CPUs)

Benchmark	#Pins	#Gates	#Arcs	OpenTimer Runtime	Our Algorithm #MDL=10		Our Algorithm #MDL=15		Our Algorithm #MDL=20	
					Runtime	Speed-up	Runtime	Speed-up	Runtime	Speed-up
leon2	4328255	1616399	7984262	2875783	4708.36	611×	5295.49ms	543×	5413.84	531×
leon3mp	3376821	1247725	6277562	1217886	5520.85	221×	7091.79ms	172×	8182.84	149×
netcard	3999174	1496719	7404006	752188	2050.60	367×	2475.90ms	304×	2484.08	303×
vga_lcd	397809	139529	756631	53204	682.94	77.9×	683.04ms	77.9×	706.16	75.3×
vga_lcd_iccad	679258	259067	1243041	66582	720.40	92.4×	754.35ms	88.3×	766.29	86.9×
b19_iccad	782914	255278	1576198	402645	2144.67	188×	2948.94ms	137×	3483.05	116×
des_perf_ispd	371587	138878	697145	24120	763.79	31.6×	766.31ms	31.5×	780.56	30.9×
edit_dist_ispd	416609	147650	799167	614043	1818.49	338×	2475.12ms	248×	2900.14	212×
mgc_edit_dist	450354	161692	852615	694014	1463.61	474×	1485.65ms	467×	1493.90	465×
mgc_matric_mult	492568	171282	948154	214980	994.67	216×	1075.90ms	200×	1113.26	193×

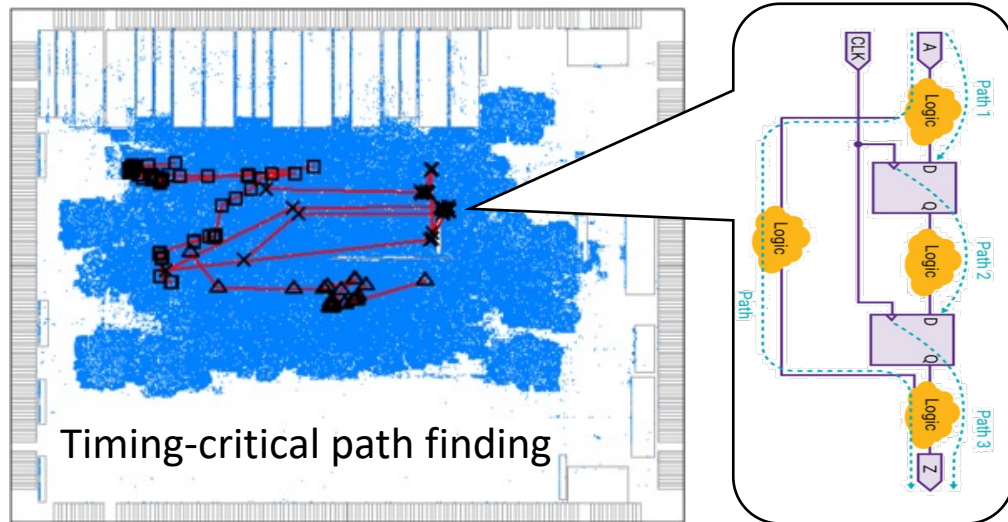
Case Study 1: Timing Analysis (cont'd)

- Comparison to existing high-performance computing systems



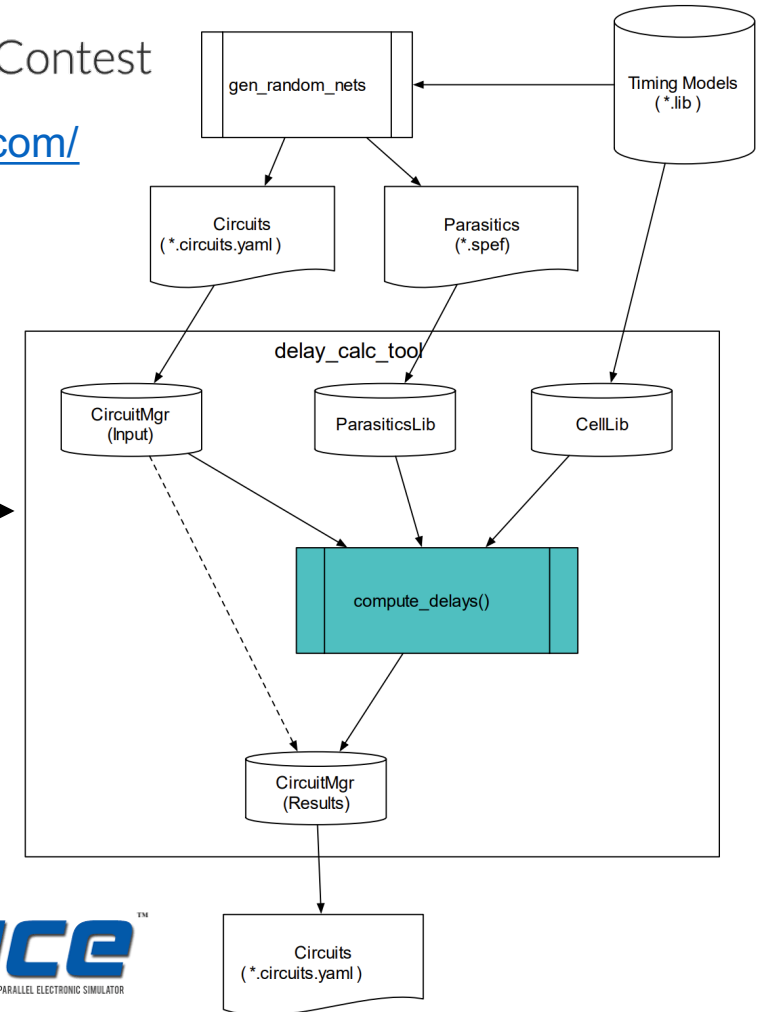
OpenTimer: Static Timing Analysis Engine

<https://github.com/OpenTimer/OpenTimer>



Tau 2021 Contest

<https://sites.google.com/view/tau-contest-2021/home>

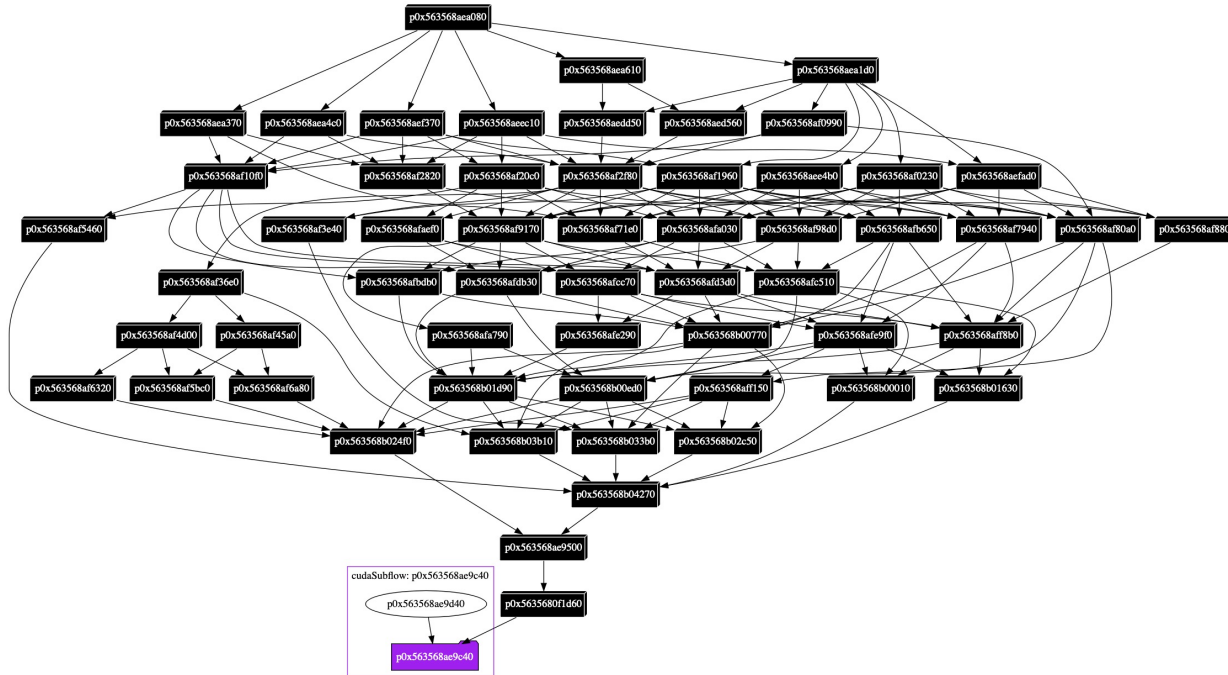


User community

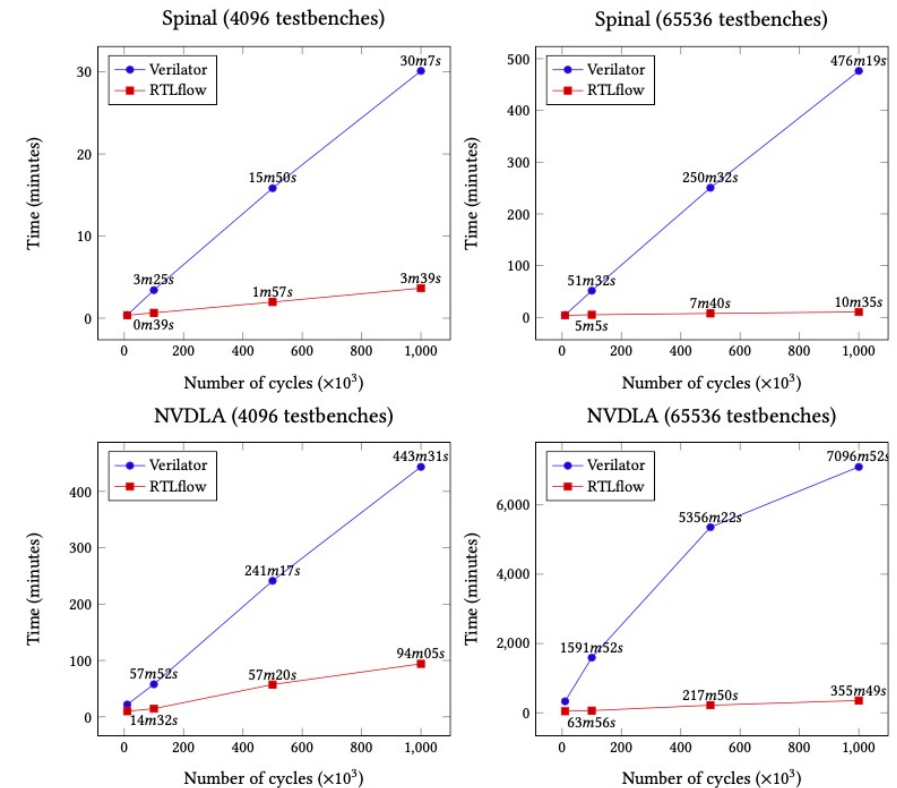


Case Study 2: RTL Simulation

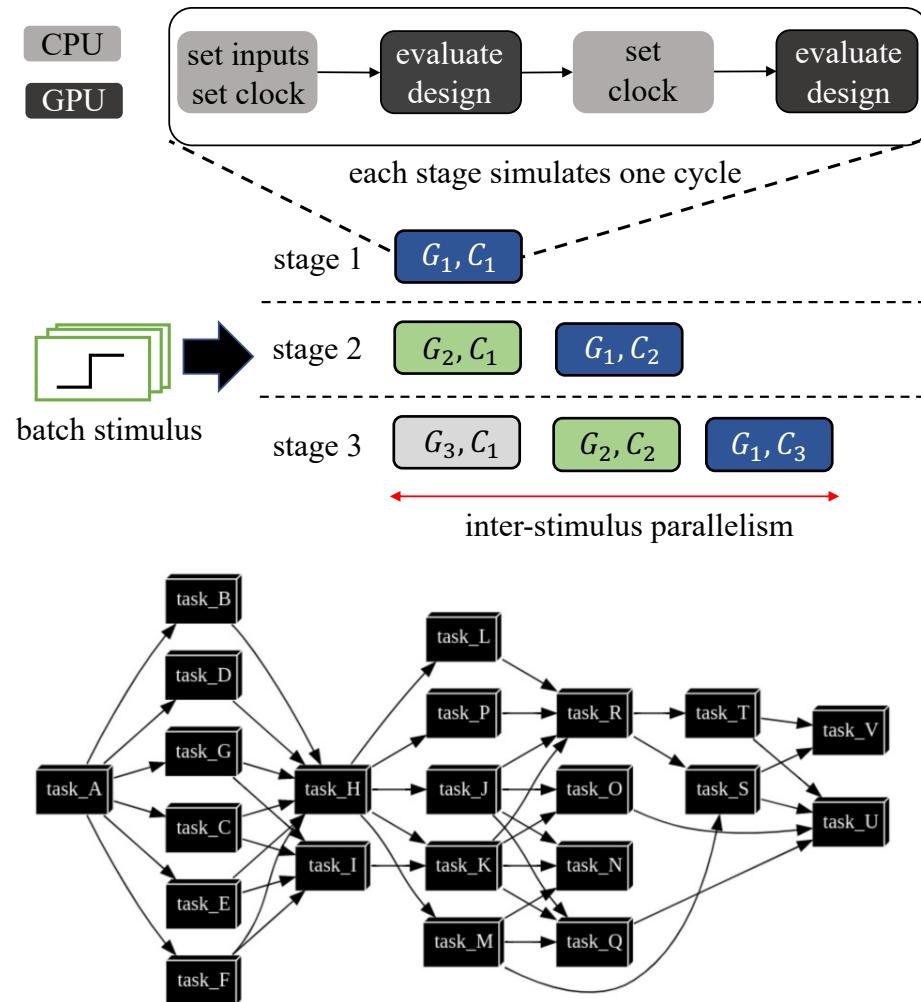
- Leverage task graph and pipeline parallelisms (i.e., RTLflow)
 - **10–500x** faster over existing RTL simulator for multiple simulation batches



Dian-Lun Lin, et al, “From RTL to CUDA: A GPU Acceleration Flow for RTL Simulation with Batch Stimulus,” *ACM ICPP*, Bordeaux, France, 2022



Case Study 2: RTL Simulation (cont'd)



#stimulus	Spinal		NVDLA	
	RTLflow ^{-p}	RTLflow	RTLflow ^{-p}	RTLflow
4096	14.7s	12.4s (↑19%)	801.2s	791.2s (↑1%)
16384	27.4s	21.4s (↑28%)	1399.2s	1098.0s (↑27%)
65536	113.8s	72.5s (↑57%)	5281.0s	2957.8s (↑79%)

Table 5: Runtime comparison in terms of improvement (↑) between RTLflow with and without pipeline scheduling (RTLflow^{-p}) for Spinal and NVDLA with 100K cycles at different numbers of stimulus.

#cycles	Spinal		NVDLA	
	stream	CUDA Graph	stream	CUDA Graph
10K	11.5s	2.3s (5×)	279.8s	106.5s (2.6×)
100K	108.0s	14.2s (7.6×)	2046.9s	791.2s (2.6×)
500K	532.9s	72.3s (7.4×)	9718.0s	3733.0s (2.6×)

Table 4: Performance advantage of CUDA Graph execution in multi-stimulus simulation workloads, measured on Spinal and NVDLA with 4096 stimulus under different numbers of cycles.

Other Industrial Applications of Taskflow

- **Quantum computing**

- Xanadu uses Taskflow in their quantum computing cloud

- **3D graphics and rendering engines**

- Methane uses Taskflow in their renderer



- **Numerical analysis**

- Deal.II uses Taskflow for advanced parallelism

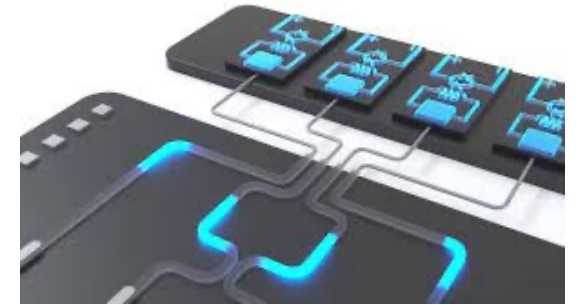
- **Computer vision**

- RevealTech uses Taskflow for real-time vision devices

- **Linear algebra**

- JetBrains uses Taskflow in their sparse matrix libraries

- ... (ME, Biochips, Imaging, FinTech, etc.)



<https://www.xanadu.ai/>



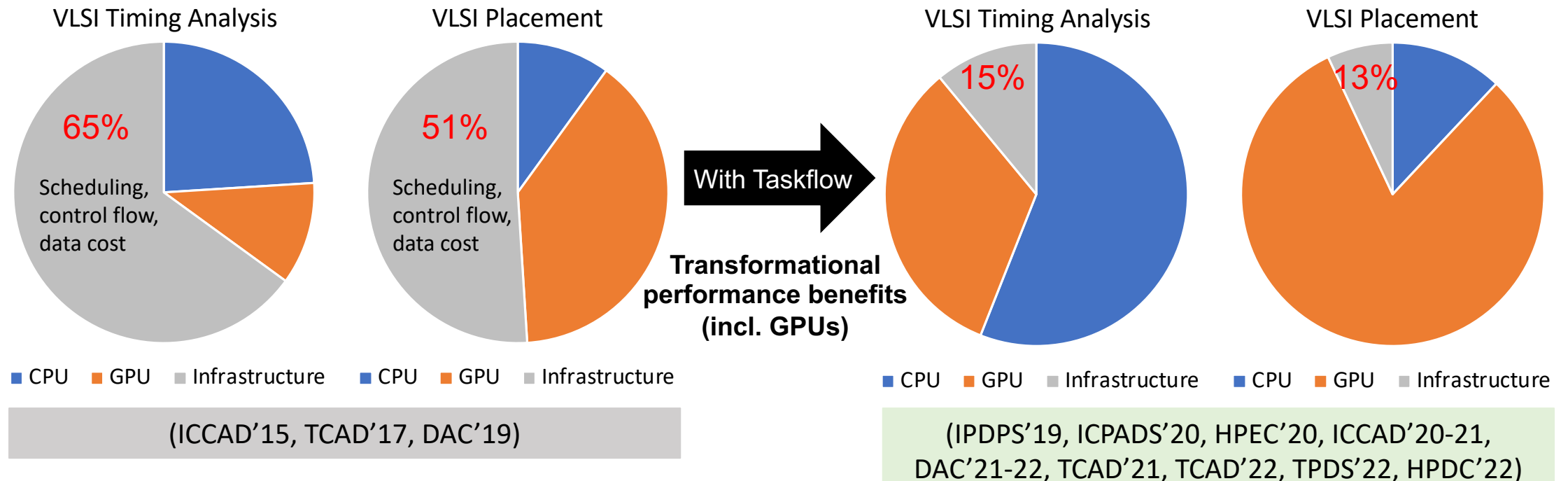
<https://www.dealii.org/>



<https://www.revealtech.ai/>

Parallel Computing Infrastructure Matters

Different models give you different implementation results. The parallel algorithm itself may run fast, but *the parallel computing infrastructure you use to implement that algorithm may dominate the entire performance.*





Use the right tool for the right job

Taskflow: <https://taskflow.github.io>

Thank You

Dr. Tsung-Wei Huang

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